Cost-efficient Virtual Optical Network Embedding for Manageable Inter-Data Center Connectivity

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Network virtualization opens the door to novel infrastructure services offering connectivity and node manageability. In this letter, we focus on the cost-efficient embedding of on-demand Virtual Optical Network (VON) requests for inter-connecting geographically distributed data centers. We present a MILP formulation that introduces flexibility in the virtual-physical node mapping in order to optimize the usage of the underlying physical resources. Illustrative results show that flexibility in the node mapping can reduce the number of add-drop ports required to serve the offered demands by 40%.

Keywords: VON, embedding, cost-efficiency, MILP.

I. Introduction

Network virtualization [1] allows network operators to capitalize on their network infrastructures by offering parts of the capacity in their physical nodes and links as infrastructure services to build virtual networks. Network virtualization fosters manageability and customizability, that is, service providers leasing a virtual network are free to configure and operate it in order to effectively deliver services to end users.

Network operators willing to strike into the virtual network service market seek to maximize profits from their network infrastructures by supporting as many virtual networks as possible, while minimizing equipment costs. With this in mind, the virtual network embedding problem [2] deals with the efficient mapping of virtual networks over a shared physical network substrate. This problem has been generally addressed in the literature. However, very few works address this problem for transparent optical networks, the main candidates to implement the future metro-core network substrates.

This work concentrates on the cost-effective embedding of Virtual Optical Networks (VONs) realizing the inter-connection of data centers (DCs) inside a cloud. We present a novel Mixed Integer Linear Programming (MILP) model to optimally embed on-demand VON requests over a high-capacity transparent optical network. The model balances wavelength utilization at physical network links together with the add-drop port utilization at physical network nodes when allocating a VON request. This serves the purpose of maximizing link resource utilization, while minimizing the number of add-drop ports to be equipped per physical node to serve the incoming VON demands.

II. VONs for data center inter-connection

In this work, we envision VONs as a promising solution for Cloud Service Providers (CSPs) lacking physical network infrastructure to inter-connect their DCs. Typically, CSPs emplace DCs in geographically distributed sites for high service reliability. In this way, they can replicate contents among them periodically, offering enhanced service availability upon failures. Besides, inter-DC workload migrations for efficiency reasons are also very frequent, which can increase data flows among DCs up to daily totals of terabytes or petabytes. In order to realize such operations, a high-capacity optical network inter-connecting DCs in a cloud is necessary. Moreover, such an optical network must provide flexibility and manageability, so as to let CSPs efficiently use it.

Thus, we believe that the VON solution perfectly suits CSPs needs. A VON is composed of virtual nodes (referred as virtual ROADMs in this paper), which are mapped over the physical nodes of the underlying network substrate (i.e., the physical
ROADMs). Virtual ROADM s are inter-connected by high-capacity virtual links, built on transparent lightpaths between the physical ROADM s over which the endpoints of the virtual links are mapped. In contrast to alternative services like bandwidth on demand or Layer-1 Virtual Private Networks (L1-VPNs), which only offer connectivity between remote sites, VONs also offer node reconfiguration capabilities. Indeed, CSPs leasing a VON can fully manage virtual nodes according to their needs, even from their own control plane. This allows them to dynamically setup/release connections between DCs at their convenience, optimizing the cloud performance.

![Figure 1. VON for data center inter-connection.](image)

In Figure 1, we depict a VON to inter-connect four DCs in a cloud. We assume that a virtual ROADM with a number of virtual add-drop ports is requested to connect each DC to the VON. A main objective of this work is to evaluate the benefits that an operator can get by introducing flexibility in the virtual-physical node mapping when embedding a VON, which should yield optimal resource utilization in both physical nodes and links. Hence, we assume that a virtual ROADM can be mapped to any physical ROADM within a geographical area around each DC (denoted as $A(VRn)$ in Figure 1, where $n$ stands for the virtual ROADM ID). The connectivity from a DC to any physical ROADM within its area can be achieved through a high-capacity access or metro network owned by the same operator provisioning the VON. Virtual links connecting virtual ROADM s are mapped on transparent lightpaths between the physical ROADM s on which the remote endpoints of the virtual links are mapped. Thus, for a virtual link, we need as many lightpaths as virtual wavelengths the virtual link requests. In this work, we assume that all virtual links of a VON request the same number of virtual wavelengths.

Two virtual-physical node mapping approaches are evaluated in this work, namely, unrestricted 1:1 and unrestricted 1:N. In the former, a virtual ROADM can only be mapped to one physical ROADM. In the latter, the virtualization middleware allows splitting a virtual ROADM over multiple physical ROADM s within its area.

### III. Related work

Even though the virtual network embedding problem has been largely addressed for electrical networks (e.g., see [2]), very few works on this topic exist for transparent optical network substrates. For instance, authors in [3] present a VON composition mechanism to serve on-demand VON requests with physical layer impairment (PLI) awareness over an all optical network substrate. In particular, fixed 1:1 virtual-physical node mappings are applied based on location requirements, and no solutions are given for performing node mappings if multiple physical nodes match the location requirements of a virtual node. Then, virtual links are mapped on the shortest paths between the physical nodes hosting the link virtual node endpoints. This work is extended to mixed-line rate networks in [4]. The work in [5] also focuses on the VON embedding problem, differentiating between transparent and opaque VONs if VON-wide optical transparency is desired or not, respectively. However, fixed 1:1 virtual-physical node mappings are also assumed, given with the VON requests. Finally, although not strictly addressing the VON embedding problem, authors in [6] focus on the energy-efficient design of integrated Virtual Infrastructures (VIs) involving both optical and IT resources, where a VI is a set of computation demands to be served in the network IT resource locations.

In contrast to the references above, the goal of this paper is to introduce flexibility in the virtual-physical node mapping, quantifying its benefits in terms of VON request blocking probability and required add-drop ports per substrate node to serve the incoming demands. We do not consider the effects of PLIs, since we assume that the distances of physical paths onto which virtual links are mapped are within the transparent reach of the optical signals (e.g., 3000 Km at 10 Gbps [7]). Nonetheless, PLI awareness could also be achieved by pruning the candidate physical paths whose end-to-end distance exceeds the transparent transmission reach. Next section details the proposed MILP formulation for optimal VON request embedding, easily adaptable to perform either unrestricted 1:1 or unrestricted 1:N node mapping approaches.

### IV. Model formulation

We model the physical substrate network as a graph $G_v(N_v, E_v)$, where $N_v$ is the set of physical nodes and $E_v$ the set of unidirectional fibers that interconnect them. The same spectral grid of $W_f$ wavelengths is assumed for each physical fiber. We also denote as $P_f$ the set of candidate paths over the physical network, namely, $p_f \in P_f$ is a candidate route to allocate a lightpath. A VON request is represented as a graph $G_v(N_v, E_v)$, where $N_v$ represents the set of virtual nodes and $E_v$ the set of
virtual links. Each virtual link of the VON requests $W_v$ wavelengths. We assume that a virtual node should be provided with as many add/drop ports as the total number of output/input virtual wavelengths from/to it. We use the terminology $a$ and $b$ to refer to the source and destination endpoints of a physical link, path or virtual link, respectively (e.g., $a(e_v)$ and $b(e_v)$ denote the s-d endpoints of $e_v$).

The decision variables of the problem are:

- $y_{n_vn_i} = \{1$ if node $n_v$ is mapped on node $n_i$; 0 otherwise $\}$, $n_v \in N_v$, $n_i \in N_i$.
- $x_{e_vp_iw} = \{1$ if $e_v$ is mapped on $p_i$ with wavelength $w$; 0 otherwise $\}$, $e_v \in E_v$, $p_i \in P_i$, $1 \leq w \leq W_i$.
- $t(n_i)$ is number of add-drop ports used in $n_i$ once the VON request is allocated.
- $t_{max}(n_v)$ is maximum add-drop port utilization among all $n_i \in A(n_v)$.
- $w(e_i)$ is number of wavelengths used at link $e_i$ once the VON request is allocated.
- $w_{max}$ is maximum wavelength usage for all $e_i \in E_i$.

The proposed MILP formulation is:

$$\min \alpha \sum_n t_{max}(n_v) + (1-\alpha)w_{max}$$

subject to:

$$\sum n_i y_{n_vn_i} = 1, \forall n_v$$  
(1a)

$$\sum n_i y_{n_vn_i} \geq 1, \forall n_v$$  
(1b)

$$y_{n_vn_i} = 0, \forall n_v, n_i \notin A(n_v)$$  
(2)

$$x_{e_vp_iw} \leq y_{a(e_v)a(p_i)}, \forall e_v, p_i, w$$  
(3)

$$x_{e_vp_iw} \leq y_{b(e_v)b(p_i)}, \forall e_v, p_i, w$$  
(4)

$$\sum_{p_i} \sum_{w} x_{e_vp_iw} = W_v, \forall e_v$$  
(5)

$$\sum_{p_i} x_{e_vp_iw} \leq \sum_{e_i} w_{max}, \forall e_i$$  
(6)

$$t_p(n_i) + \sum_{p_i} \sum_{a(p_i)=n_i} \sum_w x_{e_vp_iw} = t(n_i), \forall n_i$$  
(7)

$$t(n_i) \leq t_{max}(n_v), \forall n_v, n_i \in A(n_v)$$  
(8)

$$w_p(e_i) + \sum_{e_i} \sum_{e_i} x_{e_vp_iw} = w(e_i), \forall e_i$$  
(9)

Constraints (1) ensure that each virtual node is mapped over only one (1.a) or multiple physical nodes (1.b), depending if unrestricted 1:1 or unrestricted 1:N node mapping is applied. Constraints (2) ensure that virtual nodes are only mapped on physical nodes within their area. Constraints (3) restrict virtual link mappings to paths connecting the physical nodes over which the remote endpoints of the virtual links are mapped. Constraints (4) ensure that $W_v$ wavelengths are assigned to all virtual links. Constraints (5) are the wavelength clashing constraints. Constraints (6-9) are used to store maximum add-drop port and wavelength usages once the VON is served. $t_p(n_i)$ and $w_p(e_i)$ are the number of add-drop ports and wavelengths used in physical node $n_i$ and link $e_i$ before serving the VON, respectively.

V. Results and discussion

Using the presented model, we quantify the benefits that a network operator can achieve by introducing flexibility in the virtual-physical node mapping with unrestricted 1:1 and unrestricted 1:N node mapping approaches. Simulations are run over the extended Abilene Internet-2 network shown in Figure 2, where boxes denote areas where a DC can be located. In this network, 40 wavelengths per fiber are assumed.

To benchmark the performance of the proposed approaches, a restricted node mapping approach has been also assumed, where the physical node on which a virtual node can be mapped is fixed. We do not show the MILP formulation behind this approach due to the lack of space. Note, however, that it is considerably simpler than the one presented in section IV. Indeed, variables $y_{n_vn_i}$ and constraints (1-3) are no needed here, as all $e_v \in E_v$ have already associated the physical nodes over which $a(e_v)$ and $b(e_v)$ must be mapped. Conversely, it is here enough to enforce $x_{e_vp_iw} = 0$ for all $p_i \in P_i$ not connecting such physical nodes, which can be added as an alternative set of constraints. In any case, execution times around 1s were experienced in most occasions for all restricted and unrestricted approaches in a standard Intel Core i3 PC with 4 cores at 2.93 GHz and 2GB RAM running CPLEX v.12.2, which highlights the practicality of the proposed formulations in a relatively large transport network as the one in Figure 2.

We assume that VON requests arrive at the network in a Poisson traffic process with exponentially distributed Holding Times (HTs). VON topologies are built upon arrival as follows. For each VON request, the number of virtual ROADMs is randomly chosen from 3 to 5. Each virtual ROADM pair is connected by a bidirectional virtual link with a probability of 0.5. Once the VON topology is built, we ensure its connectivity. Otherwise, we discard it and a new topology is built. VONs request $W_v$ wavelengths for all virtual links, where $W_v$ is randomly chosen from $\{1, 2, 4\}$. Virtual nodes are assigned to the geographical areas with equiprobability, meaning that the DCs that should be connected to the VON are located in these areas. We consider that an area can allocate at most one virtual node of a VON.
Figure 3 depicts the VON blocking probability and average peak add-drop port usage per physical node for the evaluated virtual-physical node mapping approaches. Graphs are plotted as a function of the normalized VON HT, fixing the inter-arrival time of VON requests to 1 time unit. In the restricted approach, the physical node over which a virtual node must be mapped is chosen randomly among all nodes in its geographical area. In the model, $\alpha = 0.5$ has been fixed. Moreover, $P_r$ size has been limited to 6 paths for every physical node pair (i.e., the shortest ones in terms of hops).

As seen in Figure 3 (left), both unrestricted node mapping approaches yield substantially lower VON blocking for medium and high loads, as flexibility in virtual node mappings allow better resource utilization in the network. However, almost no differences are observed between them in terms of blocking probability. This is due to the fact that good wavelength usage can already be achieved by mapping virtual to physical nodes 1:1 in the geographical area. In fact, the most significant benefits from applying unrestricted 1:N are observed in the physical node add-drop port usage. Figure 3 (right) plots the average peak add-drop port usage to ensure zero blocking due to port unavailability among all network nodes. As shown, unrestricted node mapping approaches clearly lower the number of add-drop ports needed per physical node in the restricted case, which is directly translated into CAPEX savings for network operators. Among them, unrestricted 1:N interestingly decreases the number of add-drop ports needed in the restricted and unrestricted 1:1 approaches by around 40% and 20%, respectively.

VI. Conclusions

In this paper, we addressed the efficient embedding of VONs for inter-connecting geographically distributed data centers. Specifically, we investigated the benefits that a network operator can get by allowing flexibility in the virtual-physical node mapping process. We proposed two alternative flexible node mapping approaches: unrestricted 1:1 and unrestricted 1:N. The obtained results show that both approaches clearly reduce the VON request blocking and number of add-drop ports per physical node to serve the offered demands. Particularly, unrestricted 1:N reduces the number of add-drop ports per physical node by up to 40% against a restricted (i.e., fixed) node mapping, arising as an attractive solution to be adopted by future network virtualization technologies, as it can notably reduce substrate network equipment costs to network infrastructure providers.

References

OFC 2012, Mar 2011.